

Combinatorial Settlement Model: Resistance to Predators and Altruists

5th Croatian Combinatorial Days

University of Zagreb Faculty of Civil Engineering

19–20 September 2024

Mate Puljiz

University of Zagreb Faculty of Electrical Engineering and Computing

Joint work with:

Tomislav Došlić, Josip Žubrinić, Stjepan Šebek

PREDATORS AND ALTRUISTS ARRIVING ON JAMMED RIVIERA

TOMISLAV DOŠLIĆ, MATE PULJIZ, STJEPAN ŠEBEK, AND JOSIP ŽUBRINIĆ

ABSTRACT. The Riviera model is a combinatorial model for a settlement along a coastline, introduced recently by the authors. Of most interest are the so-called jammed states, where no more houses can be built without violating the condition that every house needs to have free space to at least one of its sides. In this paper, we introduce new agents (predators and altruists) that want to build houses once the settlement is already in the jammed state. Their behavior is governed by a different set of rules, and this allows them to build new houses even though the settlement is jammed. Our main focus is to detect jammed configurations that are resistant to predators, to altruists, and to both predators and altruists. We provide bivariate generating functions, and complexity functions (configurational entropies) for such jammed configurations. We also discuss this problem in the two-dimensional setting of a combinatorial settlement planning model that was also recently introduced by the authors, and of which the Riviera model is just a special case.

1. INTRODUCTION

<https://arxiv.org/pdf/2401.01225>

Joint work with



Tomislav Došlić

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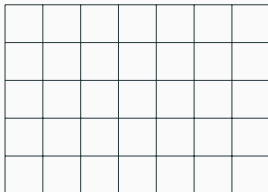
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Introduction

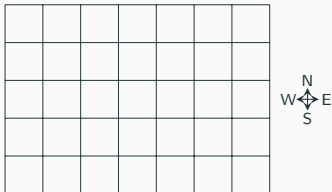
Combinatorial Settlement Model

- **A tract of land** is divided into $m \times n$ unit squares (**lots**)



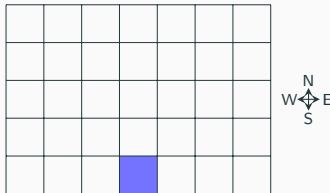
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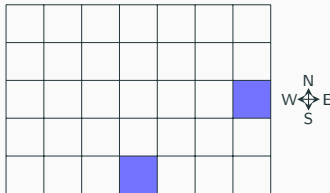
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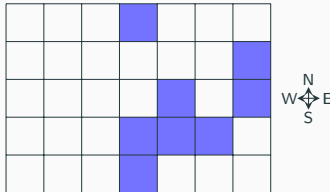
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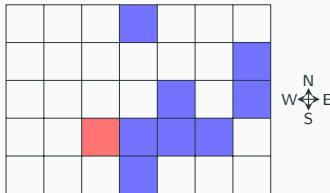
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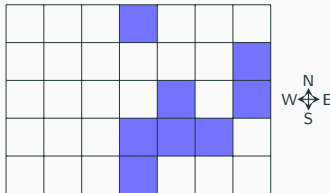
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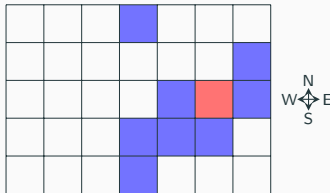
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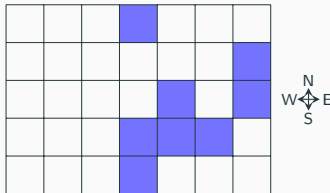
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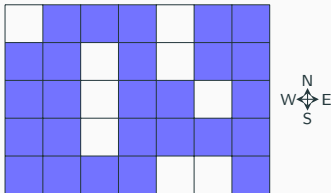
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Combinatorial Settlement Model

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- Sunlight comes from **East, South** and **West**
- Build houses ensuring that **each is adjacent to at least one unoccupied lot to its East, South, or West**
- Continue adding houses to empty lots until reaching a **jammed configuration**



Key questions

- How many jammed configurations are there?

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- What are the densities of these configurations?

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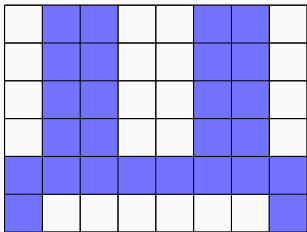
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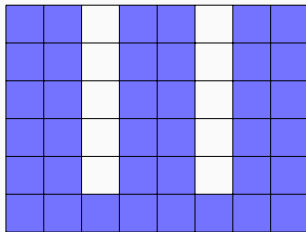
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- What is the range of possible building densities?
- Distribution over possible densities ρ ? (under which model?)
- What is the **jamming limit**, i.e., the average density of jammed configurations?
- How does the behavior change when sampling under different models?

What can be the occupancy of jammed configurations?



An inefficient configuration ($\rho \approx \frac{1}{2}$)



An efficient configuration ($\rho \approx \frac{3}{4}$)

Summary of known results ($\frac{1}{2} \leq \rho \leq \frac{3}{4}$)

Proposition (PŠŽ)

Inefficient jammed configurations on $m \times n$ grid have the following occupancy:

$$I_{m,n} = \begin{cases} \frac{mn}{2} + 2, & \text{if } n \equiv 0 \pmod{4}, \\ \frac{m(n+2)}{2}, & \text{if } n \equiv 2 \pmod{4}, \\ \frac{m(n+1)}{2} + 1, & \text{if } n \equiv 1 \pmod{2}. \end{cases}$$

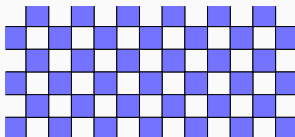
Proposition (PŠŽ)

The occupancy of efficient jammed configurations on $m \times n$ grid, $m, n \geq 2$, satisfies:

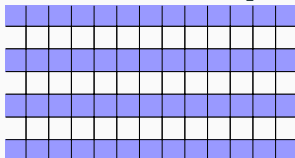
$$E_{m,n} \leq \begin{cases} mn - \lfloor \frac{n}{4} \rfloor \cdot (m-1), & \text{if } n \not\equiv 3 \pmod{4}, \\ mn - \lfloor \frac{n}{4} \rfloor \cdot (m-1) - \lfloor \frac{m}{2} \rfloor, & \text{if } n \equiv 3 \pmod{4}. \end{cases}$$

▷ P., Šebek and Žubrinić, Combinatorial settlement planning, *Contrib. Discrete Math.* **18** (2023), no. 2, 20–47

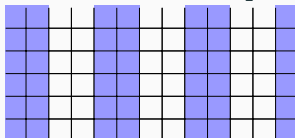
Extremal infinite grid patterns



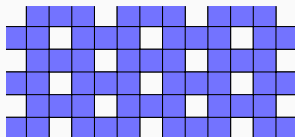
Check pattern ($\rho \approx \frac{1}{2}$)



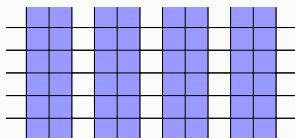
Stripe pattern ($\rho \approx \frac{1}{2}$)



Rake pattern ($\rho \approx \frac{1}{2}$)

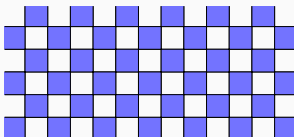


Brick pattern ($\rho \approx \frac{3}{4}$)

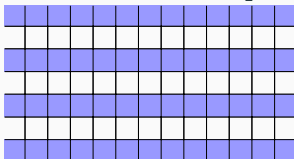


Comb pattern ($\rho \approx \frac{3}{4}$)

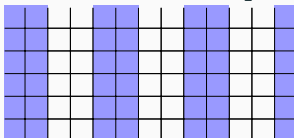
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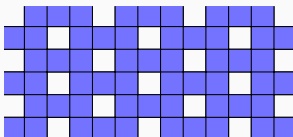
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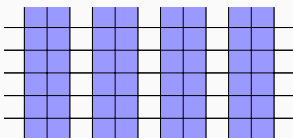
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- combine them to get densities in-between
- take a bounded piece to build a finite grid

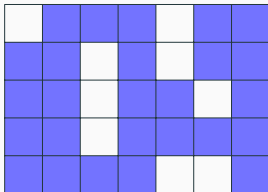
Flashback to CroCoDays 2022



Predators and altruists

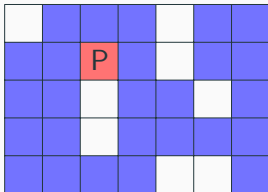
Introducing predators and altruists in the mix

- Predators and altruists are happy to build further in jammed configurations.



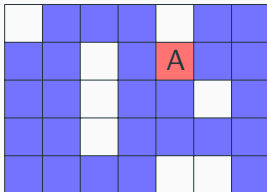
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- **Predators:** build new houses if they receive sunlight, even if it blocks others.



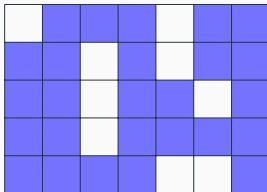
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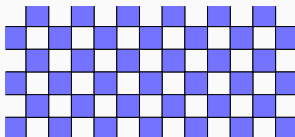


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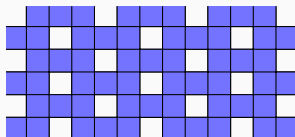
- Predators and altruists are happy to build further in jammed configurations.
- **Predators:** build new houses if they receive sunlight, even if it blocks others.
- **Altruists:** avoid blocking sunlight to others, but may build even if they don't receive sunlight.
- **Question:** How do jammed configurations which are resistant to these new asocial agents look like?



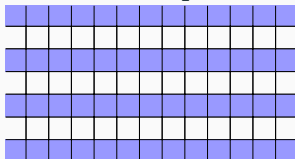
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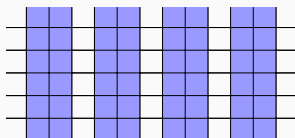
Check pattern ($\rho \approx \frac{1}{2}$) AR✗ PR✓



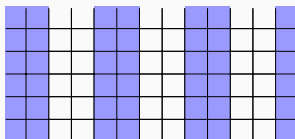
Brick pattern ($\rho \approx \frac{3}{4}$) AR✓ PR✓



Stripe pattern ($\rho \approx \frac{1}{2}$) AR✓ PR✗

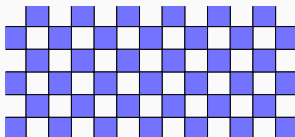


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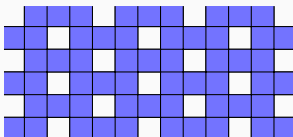


Rake pattern ($\rho \approx \frac{1}{2}$) AR✓ PR✗

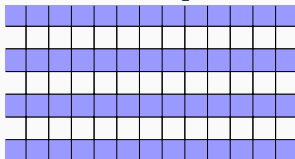
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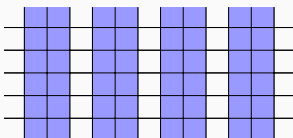
Check pattern ($\rho \approx \frac{1}{2}$) AR~~X~~ PR \checkmark



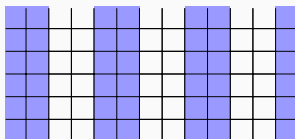
Brick pattern ($\rho \approx \frac{3}{4}$) AR \checkmark PR \checkmark



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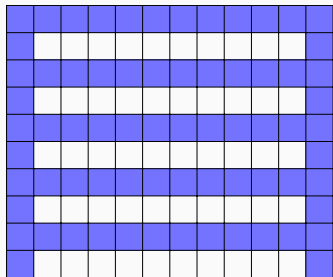


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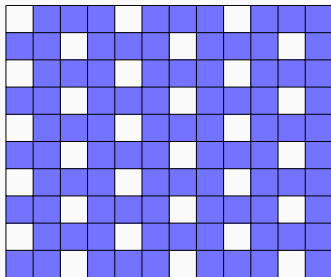
- all but Check pattern are resistant to Altruists
- just Check and Brick patterns are resistant to Predators

Jammed configurations resistant to Altruists

E.g.



Stripe pattern $\rho \approx \frac{1}{2}$

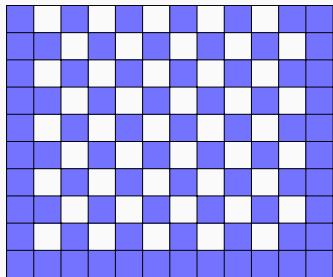


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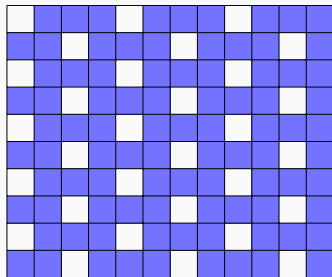
$$\frac{1}{2} \leq \rho \leq \frac{3}{4}$$

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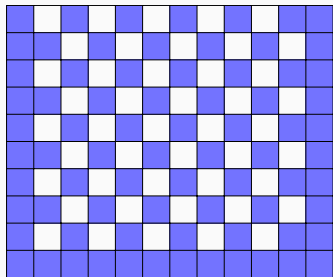


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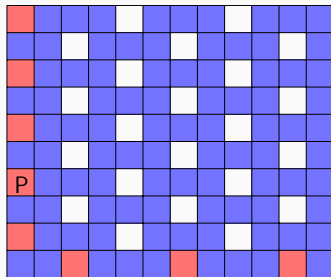
$$\frac{1}{2} \leq \rho \leq ??$$

Jammed configurations resistant to Predators

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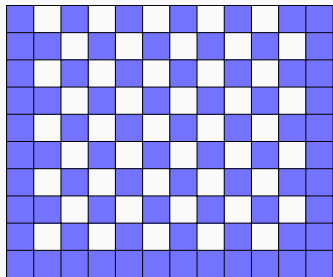


Brick pattern $\rho \approx \frac{3}{4}$ \times

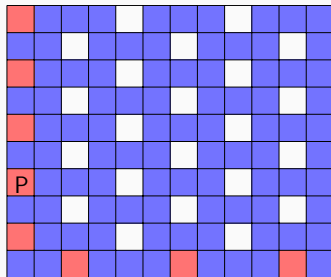
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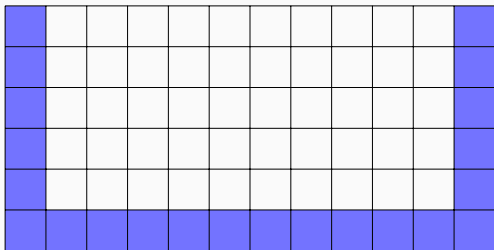


Brick pattern $\rho \approx \frac{3}{4}$ \times

$$\frac{1}{2} \leq \rho \leq \frac{2}{3}$$

Proof that $\rho \leq \frac{2}{3}$ for PR jammed configurations

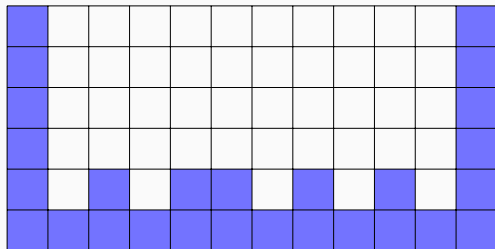
Fact 0: border must be bricked up



Proof that $\rho \leq \frac{2}{3}$ for PR jammed configurations

Fact 0: border must be bricked up

Fact 1: occupancy of the penultimate row $\leq 2 \left\lceil \frac{\# \text{ columns}}{3} \right\rceil$



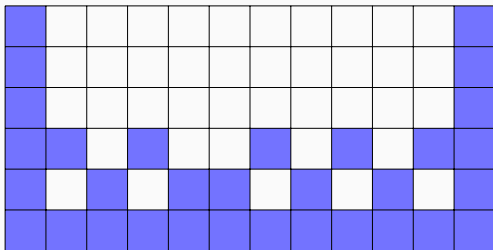
no two successive empty lots, at most two successive occupied \implies at least one empty on every three lots

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Fact 2: occupancy in successive rows (going up) can increase by at most 1



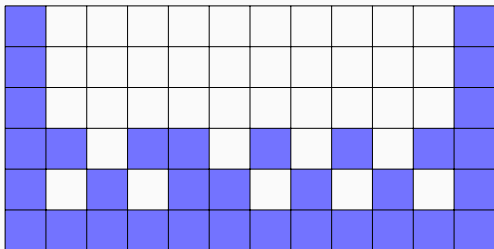
no two adjacent empty lots

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no gaps can be completely filled up \implies number of empty lots can decrease by at most one

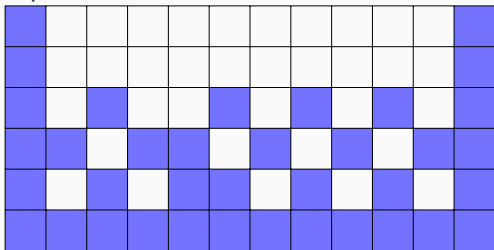
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no two adjacent empty lots

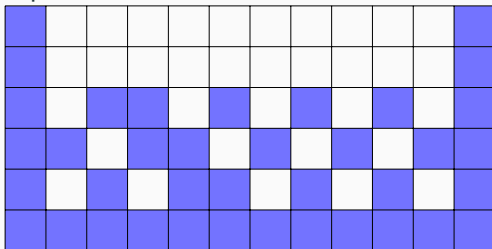
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in this case lots in columns 2 and $n - 1$ must be empty, and one more in each gap \implies number of empty lots increased by at least one

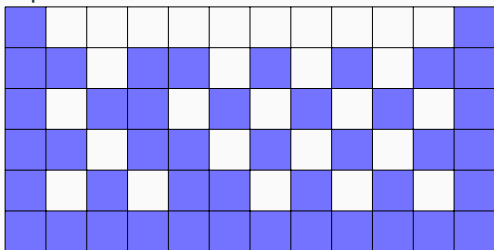
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etc.

□

Is this bound attained?

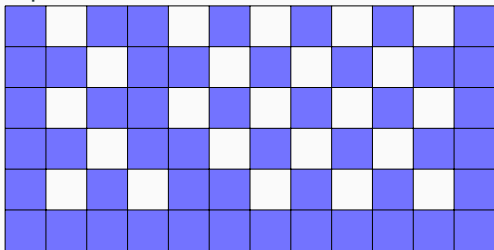
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Fact 1: occupancy of the penultimate row $\leq 2 \left\lceil \frac{\# \text{ columns}}{3} \right\rceil$

Fact 2: occupancy in successive rows (going up) can increase by at most 1

Fact 3: if it increased by 1, then it must decrease by at least 1 in the next row up



etc.

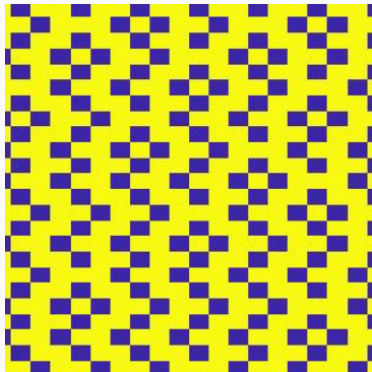
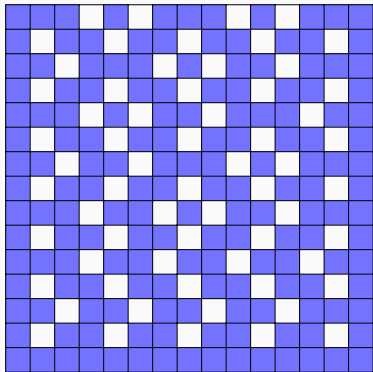
□

Is this bound attained?

Evolutionary stable configurations (ESC)

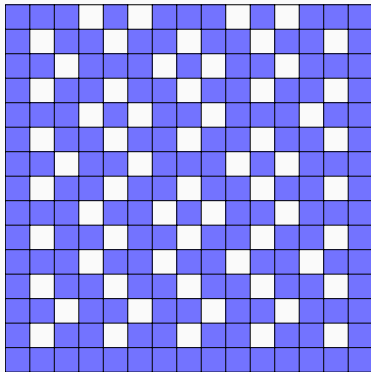
Evolutionary stable configurations – rare flowers

Evolutionary Stable Configurations: resistant to both predators and altruists.



Evolutionary stable configurations – rare flowers

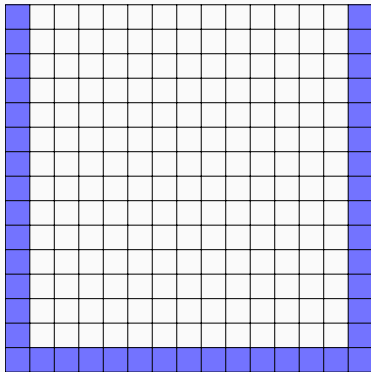
Evolutionary Stable Configurations: resistant to both predators and altruists.



Not obvious, surprising rigid structure, $\rho \approx \frac{2}{3}$, exist only for certain grid dimensions,...

Evolutionary stable configurations – rare flowers

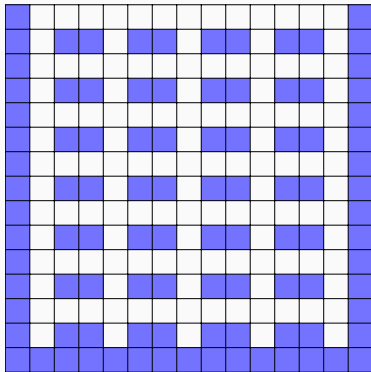
Evolutionary Stable Configurations: resistant to both predators and altruists.



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Evolutionary stable configurations – rare flowers

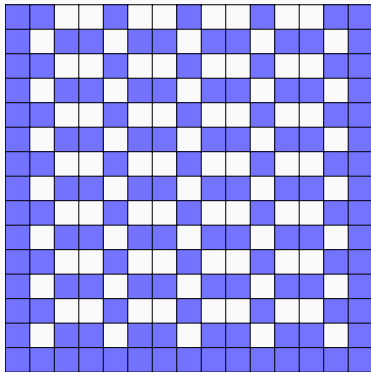
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Evolutionary stable configurations – rare flowers

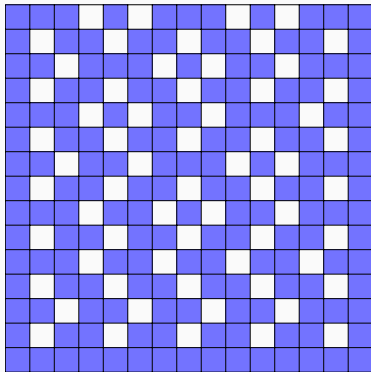
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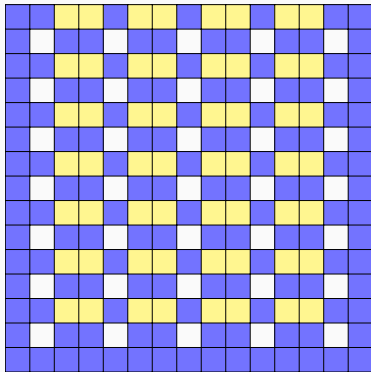
Evolutionary Stable Configurations: resistant to both predators and altruists.



Not obvious, surprising rigid structure, $\rho \approx \frac{2}{3}$, exist only for certain grid dimensions,...

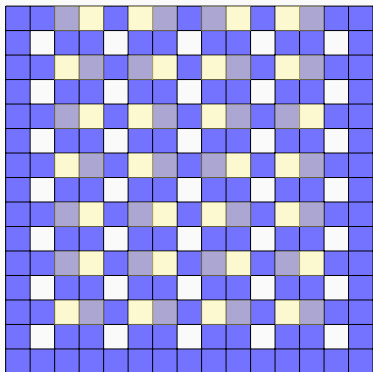
Theorem

If $m \times n$ ESC exists for $m, n > 2$, then n is divisible by 3 and m is odd, and it must have the structure below:



where exactly one in each highlighted pair of adjacent lots is occupied, and the other is empty. As a consequence, all the ES configurations have the same occupancy of $mn - \frac{(m-1)(2n-3)}{6} = \frac{2}{3}mn + \frac{1}{2}(m-1) + \frac{1}{3}n$.

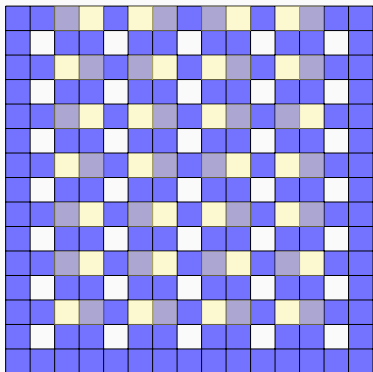
Choice is not completely arbitrary



encode
↔

R	L	R	L
L	R	L	L
R	L	L	R
L	L	R	L
R	R	L	L
R	R	R	R
L	L	L	L

Choice is not completely arbitrary



encode
 \leftrightarrow

$$\begin{bmatrix} R & L & R & L \\ L & R & L & L \\ R & L & L & R \\ L & L & R & L \\ R & R & L & L \\ R & R & R & R \\ L & L & L & L \end{bmatrix}$$

Theorem

Allowed ESCs are precisely those which (when encoded) do not contain any of the forbidden constellations:

$$\begin{array}{c} R * \\ R L \end{array}, \quad \begin{array}{c} * L \\ R L \end{array}, \quad \left| \begin{array}{c} L \\ L \end{array} \right., \quad \left| \begin{array}{c} R \\ R \end{array} \right| .$$

PREDATORS AND ALTRUISTS ARRIVING ON JAMMED RIVIERA

TOMISLAV DOŠLIĆ, MATE PULJIZ, STJEPAN ŠEBEK, AND JOSIP ŽUBRINIĆ

ABSTRACT. The Riviera model is a combinatorial model for a settlement along a coastline, introduced recently by the authors. Of most interest are the so-called jammed states, where no more houses can be built without violating the condition that every house needs to have free space to at least one of its sides. In this paper, we introduce new agents (predators and altruists) that want to build houses once the settlement is already in the jammed state. Their behavior is governed by a different set of rules, and this allows them to build new houses even though the settlement is jammed. Our main focus is to detect jammed configurations that are resistant to predators, to altruists, and to both predators and altruists. We provide bivariate generating functions, and complexity functions (configurational entropies) for such jammed configurations. We also discuss this problem in the two-dimensional setting of a combinatorial settlement planning model that was also recently introduced by the authors, and of which the Riviera model is just a special case.

1. INTRODUCTION

<https://arxiv.org/pdf/2401.01225>

Thank You!

Forbidden constellations

$$\begin{array}{cccc|cccc|...|cccc|cccc} 0 & * & * & * & * & * & * & * & \dots & * & * & * & * & * & * & * \\ * & \underline{0} & * & * & * & * & * & * & \dots & * & * & * & 1 & 1 & 1 \\ * & * & \underline{0} & 1 & 1 & 0 & 1 & \dots & 1 & 0 & 1 & 1 & 0 & 1 \end{array}$$

(a) Type I

$$\begin{array}{cccc|cccc|...|cccc|cccc} * & \underline{0} & * & * & * & * & * & * & \dots & * & * & * & 1 & 1 & 1 \\ 0 & * & \underline{0} & 1 & 1 & 0 & 1 & \dots & 1 & 0 & 1 & 1 & 0 & 1 \end{array}$$

(b) Type II

$$\begin{array}{cccccccc} 0 & * & * & * & * & * & * & \dots \\ * & \underline{0} & * & * & * & * & * & \dots \\ * & * & \underline{0} & * & * & * & * & \dots \\ * & * & 1 & 1 & 1 & 1 & 1 & \dots \end{array}$$

(c) Type III

$$\begin{array}{cccccccccccc} 0 & * & * & * & * & * & * & * & * & * & * & \dots \\ * & \underline{0} & * & * & * & * & * & * & * & * & * & \dots \\ * & * & \underline{0} & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & \dots \end{array}$$

(d) Type IV

$$\begin{array}{cccccccc} \dots & * & * & * & * & \underline{1} & * & * & * & * & \dots \\ \dots & * & * & * & * & * & * & * & * & * & \dots \\ \dots & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & \dots \end{array}$$

(e) Type V

Forbidden constellations

$$\begin{array}{cccc|cccc|...|cccc|cccc}
 0 & * & * & * & * & * & * & * & \dots & * & * & * & * & * & * & * \\
 * & \underline{0} & * & * & * & * & * & * & \dots & * & * & * & * & * & * & * \\
 * & * & \underline{0} & 1 & 1 & 0 & 1 & \dots & 1 & 0 & 1 & \dots & 1 & 0 & 1 & 1 & 0 & 1
 \end{array}$$

(a) Type I

$$\begin{array}{cccc|cccc|...|cccc|cccc}
 * & \underline{0} & * & * & * & * & * & * & \dots & * & * & * & * & * & * & 1 & 1 & 1 \\
 0 & * & \underline{0} & 1 & 1 & 0 & 1 & \dots & 1 & 0 & 1 & \dots & 1 & 0 & 1 & 1 & 0 & 1
 \end{array}$$

(b) Type II

$$\begin{array}{cccccccc}
 0 & * & * & * & * & * & * & \dots \\
 * & \underline{0} & * & * & * & * & * & \dots \\
 * & * & \underline{0} & * & * & * & * & \dots \\
 * & * & 1 & 1 & 1 & 1 & 1 & \dots
 \end{array}$$

(c) Type III

$$\begin{array}{cccccccccccc}
 0 & * & * & * & * & * & * & * & * & * & * & \dots \\
 * & \underline{0} & * & * & * & * & * & * & * & * & * & \dots \\
 * & * & \underline{0} & 1 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & \dots
 \end{array}$$

(d) Type IV

$$\begin{array}{cccccccc}
 \dots & * & * & * & * & \underline{1} & * & * & * & * & \dots \\
 \dots & * & * & * & * & * & * & * & * & * & \dots \\
 \dots & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 1 & \dots
 \end{array}$$

(e) Type V

... and additionally East↔West mirrored versions of Type I–IV

Riviera model (1D)

