

# A characterisation of compatible state space aggregations

Mate Puljiz  
University of Birmingham

joint work with  
Chris Good, David Parker, and Jonathan E. Rowe

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## What are we studying?

- ▶ When can we *simplify* an iterated system

$$\mathbf{x}_{n+1} = f(\mathbf{x}_n), \text{ where } f : \mathbb{R}^N \rightarrow \mathbb{R}^N?$$

- ▶ E.g.

$$x_{n+1} = x_n^2,$$

$$y_{n+1} = y_n^2 + 2x_n y_n,$$

$$z_{n+1} = z_n + w_n,$$

$$w_{n+1} = 0,$$

can be simplified to

$$X_{n+1} = X_n^2,$$

$$Y_{n+1} = Y_n,$$

where  $X_i = x_i + y_i$  and  $Y_i = z_i + w_i$ .

# Why are we studying this?

- ▶ Very general framework
- ▶ Simplifying real-life models (model reduction)
- ▶ Discovering underlying regularities of a model
- ▶ Efficient usage of computational resources

# Our goal

- ▶ Identify a class of iterated systems of interest
- ▶ Interpret these as a parametrised set of models for artificial chemistries
- ▶ Given one such model, characterise its compatible aggregations in terms of the intrinsic parameters of the model

# Notation

- ▶ Aggregation – a partition of the set of variables (sometimes the induced equivalence relation on  $\{1, 2, \dots, N\}$ )
  - ▶  $\{\{x, y\}, \{z, w\}\}$  above
  - ▶ Can be encoded by a 0-1 column stochastic matrix
  - ▶ E.g.  $\Xi = \begin{bmatrix} 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}$  above as
$$(X_i, Y_i)^T = \Xi \cdot (x_i, y_i, z_i, w_i)^T$$
- ▶ Compatible aggregation – an aggregation for which the coarse grained description is well-defined

# Class 1 – Markov Chains

## Parameter set

$M$  – an  $N \times N$  column stochastic matrix

## Iterated system

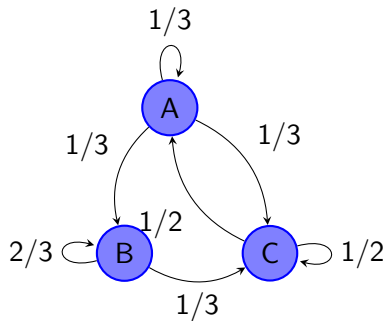
$$\mathbf{p}_{n+1} = M\mathbf{p}_n,$$

where  $\mathbf{p} = (p^1, p^2, \dots, p^N)$

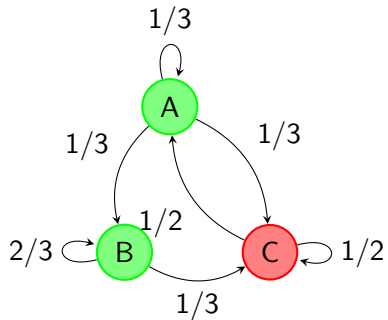
## Proposition (Folklore, see Kemeny & Snell [9])

*An equivalence relation  $\equiv$  on the set of states  $\{1, 2, \dots, N\}$  is a compatible aggregation if and only if any two equivalent states have exactly the same outgoing transition probabilities towards any of the blocks of the (induced) partition.*

# A picture

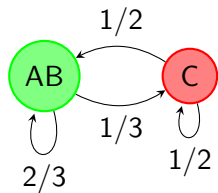
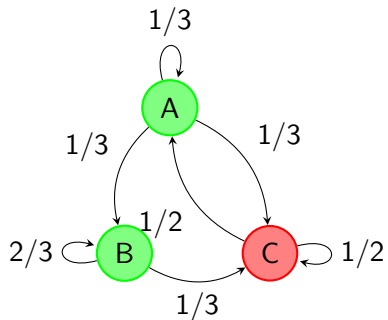


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# Class 1 – Markov Chains

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Note that if we denote by  $\Xi$  the associated 0-1 matrix, the condition above could be equally written as

$$\Xi(M_i) = \Xi(M_j) \quad \text{whenever} \quad i, j \in \{1, 2, \dots, N\} \text{ such that } i \equiv j,$$

where  $M_k$  is the  $k^{\text{th}}$  column of  $M$ .

# Class 1 – Markov Chains

More abstractly

$\Xi(\alpha_v) = \Xi(\alpha_w)$  whenever  $v, w \in \mathbb{Z}_+^N$  such that  $\Xi(v) = \Xi(w)$ ,

where  $\alpha_v = \partial_v M(0, \dots, 0) = \begin{cases} M_i, & \text{if } v = e_i, \quad (\partial_{e_i} = \partial_i) \\ (0, \dots, 0)^T, & \text{otherwise.} \end{cases}$

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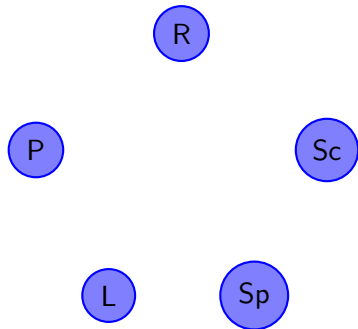
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## Class 2 – Binary Tournaments

### Parameter set

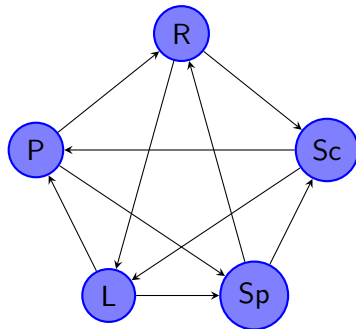
$\leftarrow$  – an orientation on a complete graph on  $N$  vertices (a finite set of parameters)



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## Iterated system

$$p_{n+1}^i = (p_n^i)^2 + 2 \sum_{i \leftarrow k} p_n^k \cdot p_n^i = (T(p_n^1, \dots, p_n^N))_i,$$

for  $i \in \{1, 2, \dots, N\}$

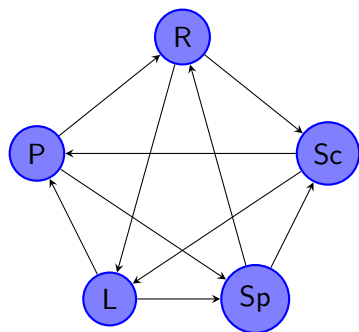
## Class 2 – Binary Tournaments

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*An equivalence relation  $\equiv$  on  $\{1, 2, \dots, N\}$  is a compatible aggregation for the iterated system above if and only if*

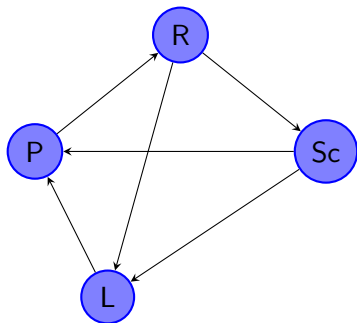
$$i \equiv j \equiv k \quad \text{whenever} \quad i \equiv k \text{ and } i \leftarrow j \leftarrow k.$$

# A picture

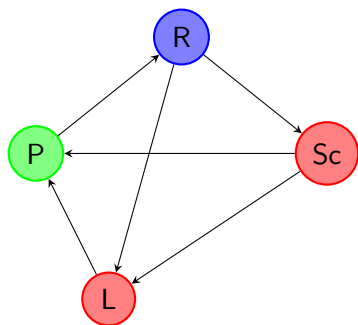




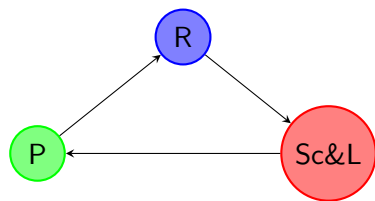
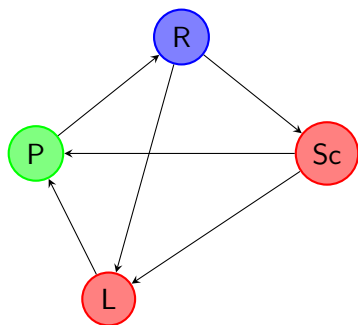
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$$\Xi(\alpha_v) = \Xi(\alpha_w) \quad \text{whenever} \quad v, w \in \mathbb{Z}_+^N \text{ such that } \Xi(v) = \Xi(w),$$

where

$$\alpha_v = \partial_v T(0, \dots, 0) = \begin{cases} 2e_i, & \text{if } v = 2e_i, & (\partial_{2e_i} = \partial_{ii}) \\ 2e_i, & \text{if } v = e_i + e_k \text{ and } i \leftarrow k, \\ (0, \dots, 0)^T, & \text{otherwise.} & (\partial_{e_i+e_k} = \partial_{ik}) \end{cases}$$

## Class 2 v2 – 2<sup>nd</sup> order Chemistry

### Parameter set

$\alpha_v$  – 2× the vector of probabilities of getting a certain product in a reaction represented by  $v$ , where  $v \in \mathbb{Z}_+^N$  s.t.  $|v| = 2$

### Iterated system

$$\mathbf{p}_{n+1} = T(\mathbf{p}_n),$$

where  $\mathbf{p} = (p^1, p^2, \dots, p^N)$  and

$$T(\mathbf{p}) = \sum_{v \in \mathbb{Z}_+^N: |v|=2} \frac{1}{v!} \cdot \alpha_v \cdot \mathbf{p}^v$$

## Class 2 v2 – 2<sup>nd</sup> order Chemistry

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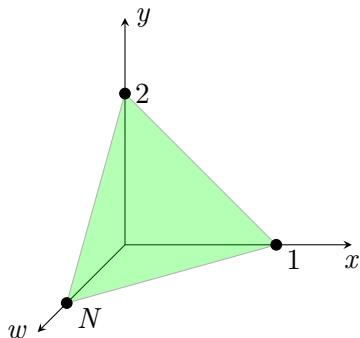
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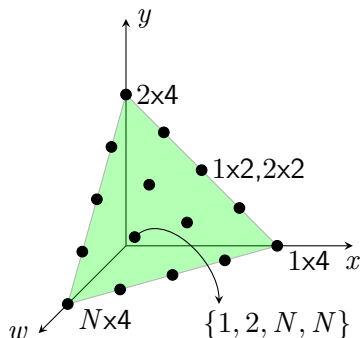
## Discretisation and Modelling issues

- ▶  $\Omega$  —  $\{1, 2, \dots, N\}$
- ▶  $T$  — dynamics on  $\Delta_N = M(\Omega)$
- ▶  $R$  — size of a total population in simulation



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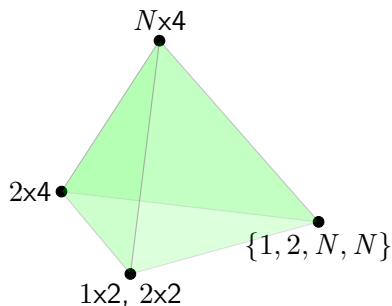
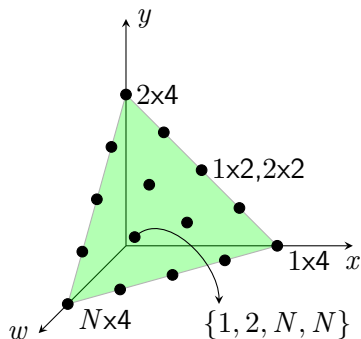
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- ▶  $\frac{1}{R} X_N^R \subset \Delta_N$  — set of finite populations
- ▶  $M(\frac{1}{R} X_N^R) = \Delta_{C(N,R)}$  — set of probabilities over finite populations



## Different models

A countable family of independent random variables

$$\left\{ Z_{\alpha}^{k,i}, Y_{\beta}^{k,i} : k \in \mathbb{N}, i \in \mathbb{N}, \alpha \in T\left(\frac{1}{R}X_N^R\right), \beta \in \frac{1}{R}X_N^R \right\},$$

where

$$Z_{\alpha}^{k,i} \sim \begin{pmatrix} 1 & 2 & \dots & N \\ \alpha_1 & \alpha_2 & \dots & \alpha_N \end{pmatrix} \text{ and } Y_{\beta}^{k,i} \sim \begin{pmatrix} 1 & 2 & \dots & N \\ \beta_1 & \beta_2 & \dots & \beta_N \end{pmatrix}$$

$$\frac{1}{R}V_k = F\left(\frac{1}{R}V_{k-1}; Z_{T(\frac{1}{R}V_{k-1})}^{k,1}, Z_{T(\frac{1}{R}V_{k-1})}^{k,2}, \dots; Y_{\frac{1}{R}V_{k-1}}^{k,1}, Y_{\frac{1}{R}V_{k-1}}^{k,2}, \dots\right),$$

for some  $F: \frac{1}{R}X_N^R \times \Omega^{\infty} \times \Omega^{\infty} \rightarrow \frac{1}{R}X_N^R$  which is assumed to be compatible with  $\Xi_{\pi} \times \pi^{\infty} \times \pi^{\infty}$  (for any quotient map  $\pi: \Omega \rightarrow \tilde{\Omega}$ )

# Different models

Generational model

$$F_2 \left( \frac{1}{R} \mathbf{v}; i_1, i_2, \dots; j_1, j_2, \dots \right) = \frac{1}{R} (\mathbf{e}_{i_1} + \mathbf{e}_{i_2} + \dots + \mathbf{e}_{i_R})$$

One at a time reaction

$$F_1 \left( \frac{1}{R} \mathbf{v}; i_1, i_2, \dots; j_1, j_2, \dots \right) = \frac{1}{R} (\mathbf{v} + \mathbf{e}_{i_1} - \mathbf{e}_{j_1})$$

## Proposition

*If  $\Xi$  is a compatible aggregation for  $T$ , then the Markov chains induced by the modelling functions  $F_1$  (Vose [20]) or  $F_2$  (or other compatible  $F$ ) can be aggregated in the same way.*

## Class 3 – A very general heuristic model

### Parameter set

$\alpha_v$  –  $2 \times$  the vector of probabilities of getting a certain product in a reaction represented by  $v$ , where  $v \in \mathbb{Z}_+^N$  s.t.  $|v| = 2$

### Iterated system

$$\mathbf{p}_{n+1} = T(\mathbf{p}_n),$$

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## Class 3 – Compatible aggregations

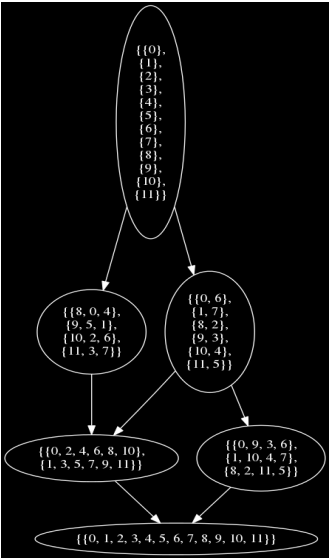
### Theorem

*An equivalence relation  $\equiv$  on  $\{1, 2, \dots, N\}$  is a compatible aggregation for the iterated system above if and only if*

$$\Xi(\alpha_v) = \Xi(\alpha_w) \quad \text{whenever} \quad v, w \in \mathbb{Z}_+^N \text{ such that } \Xi(v) = \Xi(w).$$

Note that  $\alpha_v = \partial_v T(0, \dots, 0)$ .

# Modulo 12 Game

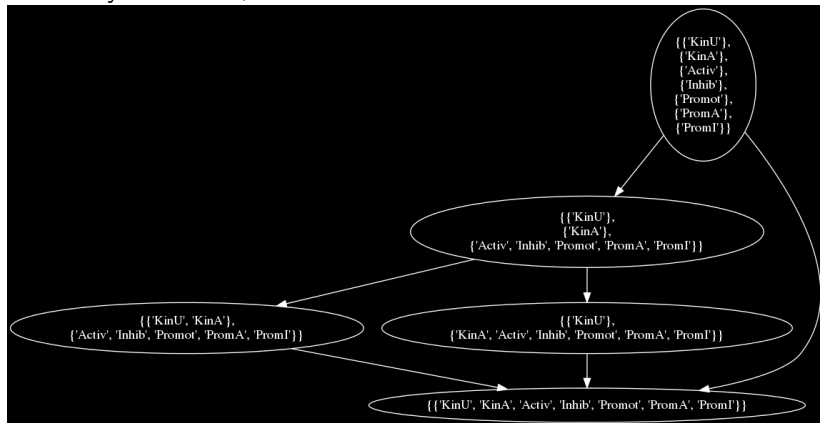


$$x + y \rightarrow (x +_{12} y)$$

$$\Omega = \{0, 1, \dots, 11\}$$

# An application to Kinetochore

Model by R. Henze, and B. Ibrahim





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Thank you for your attention!

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